## Programming the Demirci-Selçuk Meet-in-the-Middle Attack with Constraints

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## Outlines

(1) Introduction
(2) Modelling the MITM attack

3 MITM and Impossible differential application in design
4. Conclusion

## Outline

(1) Introduction

- Searching methods
- Distinguisher of Demirci-Selçuk MITM
- Key recovery attack of MITM
(2) Modelling the MITM attack

3) MITM and Impossible differential application in design
(4) Conclusion

## Automatic Cryptanalysis

- Dedicated search
- MILP,CP,SAT,SMT


## Searching methods for MITM

- Demirci-Selçuk MITM, FSE 2008.
- Derbez and Fouque: Dedicated search algorithm
- Li Lin, Wenling Wu: General model based on MILP


## MITM Distinguisher



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$$

$\Delta_{\mathrm{E}}(A, B):\left\{C^{0}[B] \oplus C^{1}[B], C^{0}[B] \oplus C^{2}[B], \ldots, C^{0}[B] \oplus C^{N-1}[B]\right\}$


- Random Cipher: $\mathcal{N}_{R}$

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Distinguisher: $\left(A, B, \mathcal{N}_{E}\right)$

## Structure of the attack

- a cipher is divided in three keyed permutations: $E_{0}, E_{1}, E_{2}$
- Construct distinguisher $\left(A, B, \mathcal{N}_{E}\right)$ at $E_{1}$



## Outline

(1) Introduction
(2) Modelling the MITM attack

- Modelling the distinguisher
- Modelling the Key-Recovery Process

3) MITM and Impossible differential application in design
(4) Conclusion

## Variables



- $\operatorname{Var}(\mathrm{X})$ describe the forward differential
- $\operatorname{Var}(\mathrm{Y})$ describe the backward determination
- $\operatorname{Var}(Z)$ models the relation between $\operatorname{Var}(X)$ and $\operatorname{Var}(Y)$


## Forward differential

## Variables Var (X)

$X_{r}[j]=0$ iff $P_{r}^{0}[j] \oplus P_{r}^{i}[j]=0, \forall i \in 1, \ldots, N-1$.


$$
\begin{aligned}
x_{2} & =x_{0} \\
2 x_{3} & \geq x_{0}+x_{1} \\
x_{3} & \leq x_{0}+x_{1}
\end{aligned}
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## Backward determination

## Variables Var (Y)



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\begin{aligned}
y_{2}+y_{3} & \leq 2 y_{0} \\
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## Constraints for Var (Z)

Variables $\operatorname{Var}(\mathrm{Z})$ describe the relation between $\operatorname{Var}(\mathrm{X})$ and $\operatorname{Var}(\mathrm{Y})$ :

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Z_{r}[j]=1 \text { iff } X_{r}[j]=Y_{r}[j]=1
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objective function: Minimize $\sum_{r=r_{0}+1}^{r_{0}+r_{1}-1} Z_{2 r}$

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Round 1
Round 2


$$
M C=\left(\begin{array}{llll}
1 & 0 & 1 & 1 \\
1 & 0 & 0 & 0 \\
0 & 1 & 1 & 0 \\
1 & 0 & 1 & 0
\end{array}\right)
$$



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Table 2: An enumeration of all $\mathcal{D S}$-MITM distinguishers for 10.5 -round SKINNY-
128-384 with $40 \leq \operatorname{Deg}(\mathcal{A}, \mathcal{B}) \leq 48$.

| No. | $\mathcal{A}$ | $\mathcal{B}$ | $\operatorname{Deg}(\mathcal{A}, \mathcal{B})$ | No. | $\mathcal{A}$ | $\mathcal{B}$ | $\operatorname{Deg}(\mathcal{A}, \mathcal{B})$ | No. | $\mathcal{A}$ | $\mathcal{B}$ | $\operatorname{Deg}(\mathcal{A}, \mathcal{B})$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | $[15][4]$ | 40 | 21 | $[13][6,4]$ | 45 | 41 | $[13]$ | $[5]$ | 46 |  |  |
| 2 | $[12][5]$ | 40 | 22 | $[14]$ | $[7,5]$ | 45 | 42 | $[12]$ | $[4]$ | 46 |  |
| 3 | $[13][6]$ | 40 | 23 | $[13]$ | $[6,4]$ | 45 | 43 | $[14]$ | $[6]$ | 46 |  |
| 4 | $[14][7]$ | 40 | 24 | $[15][4,6]$ | 45 | 44 | $[15]$ | $[7]$ | 46 |  |  |
| 5 | $[15][5]$ | 42 | 25 | $[13]$ | $[5]$ | 45 | 51 | $[13]$ | $[4,6]$ | 47 |  |
| 6 | $[12][6]$ | 42 | 26 | $[15]$ | $[6]$ | 45 | 52 | $[12]$ | $[7,5]$ | 47 |  |
| 7 | $[13][7]$ | 42 | 27 | $[14]$ | $[4]$ | 45 | 53 | $[14]$ | $[5,7]$ | 47 |  |
| 8 | $[14][4]$ | 42 | 28 | $[13]$ | $[4]$ | 45 | 54 | $[15]$ | $[6,4]$ | 47 |  |
| 9 | $[13][5]$ | 43 | 29 | $[14]$ | $[5]$ | 45 | 49 | $[13]$ | $[6]$ | 47 |  |
| 10 | $[14][6]$ | 43 | 30 | $[14]$ | $[6]$ | 45 | 50 | $[13]$ | $[6]$ | 47 |  |
| 11 | $[12][4]$ | 43 | 31 | $[12]$ | $[4]$ | 45 | 51 | $[14]$ | $[7]$ | 47 |  |
| 12 | $[15][7]$ | 43 | 32 | $[15]$ | $[5]$ | 45 | 52 | $[12]$ | $[5]$ | 47 |  |
| 13 | $[12][7]$ | 44 | 33 | $[13]$ | $[7]$ | 45 | 53 | $[12]$ | $[5]$ | 47 |  |
| 14 | $[13][4]$ | 44 | 34 | $[12]$ | $[6]$ | 45 | 54 | $[14]$ | $[7]$ | 47 |  |
| 15 | $[12][7]$ | 44 | 35 | $[15]$ | $[7]$ | 45 | 55 | $[15]$ | $[4]$ | 47 |  |
| 16 | $[13][4]$ | 44 | 36 | $[12]$ | $[7]$ | 45 | 56 | $[15]$ | $[4]$ | 47 |  |
| 17 | $[13][4]$ | 44 | 37 | $[14]$ | $[4,6]$ | 46 | 57 | $[15]$ | $[7,5]$ | 48 |  |
| 18 | $[14][5]$ | 44 | 38 | $[13][7,5]$ | 46 | 58 | $[14]$ | $[6,4]$ | 48 |  |  |
| 19 | $[14][5]$ | 44 | 39 | $[15][5,7]$ | 46 | 59 | $[12]$ | $[4,6]$ | 48 |  |  |
| 20 | $[13][4]$ | 44 | 40 | $[12]$ | $[6,4]$ | 46 | 60 | $[13]][5,7]$ | 48 |  |  |

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## New 0-1 variables Var(M) and Var(W)

Round 0

$M C^{-1}=\left(\begin{array}{llll}0 & 1 & 0 & 0 \\ 0 & 1 & 1 & 1 \\ 0 & 1 & 0 & 1 \\ 1 & 0 & 0 & 1\end{array}\right)$
$\operatorname{Var}(W)$ : Forward determination


Var(M): Backward differential

| Round 2 |  | SB,AC |  |  |  | Round3 |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  |  |  |  |  |  |  |  |
|  |  |  |  |  |  | MC |  |  |  |
|  |  | AK, SR |  |  |  |  |  |  |  |
|  |  |  |  |  |  |  |  |  |  |



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$\operatorname{Var}(W)$ : Forward determination



Distinguisher


11-round Distinguisher

## 









1-round Distinguisher

## Key bridging technique



## Key bridging technique

${ }^{0}{ }^{1}{ }^{2}{ }^{2}{ }^{3}{ }^{4} \frac{5}{5}$


Round 3


 Shift to the left by 29-bit

Round



the left by 29-bit
Round
 |1| | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |

## Outline

(1) Introduction
(2) Modelling the MITM attack
(3) MITM and Impossible differential application in design

- Results of Lblock
- Results of TWINE
(4) Conclusion


## LBlock

- LBlock

- $8!=40320$ variants ciphers against MITM and ID


## LBlock

- LBlock

- $8!=40320$ variants ciphers against MITM and ID


## Results of LBlock

All exist 14-round ID distinguisher
32 permutations are good:

- no 15 -round ID distinguisher
- strong against the MITM Distinguisher


## TWINE

## TWINE Cipher:



Enumeration: $22 \cdot 8$ !

## TWINE

## TWINE Cipher:



Enumeration: $22 \cdot 8$ !

## Results of TWINE

- 144 permutations: no 15-round ID Distinguisher.
- 84 permutations are good in the view of MITM.
- 12 permutations are best: no 11-round MITM distinguisher


## Outline

## (2) Modelling the MITM attack

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## Conclusion

Conclusion

- modelling the MITM attack
- ID and MITM for variants cipher of LBlock and TWINE

Future Work - Differential enumaraion

- Key Bridging


## Thanks for your attention.

